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**Data Encryption and Decryption using Modified RSA Cryptography Based on Multiple
Public Keys and 'n' Prime Number**

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Abstract

Now a day, we are having a great dependence on computer and network for communication. The security of computer communication is related to the whole world and everybody. Cryptography is the art and science of achieving security by encoding message to make them non-readable to secure data or information transmits over the network. In this paper introduced modified RSA approach based on multiple public keys and n prime number. RSA algorithm is mostly used in the popular implementation of public key cryptography. In public key cryptography two different keys are generated in RSA one keys is used in encryption data and other corresponding key used for decryption. No other key decrypt the data. Even if it is efficient algorithm it is vulnerable to other person. With the help of all brute force attacks can obtain private keys. In this research paper new approach we used n prime number and multiple public keys. Which is not easily crack able. In here implementation RSA algorithm. Using some mathematical logic integer factorization and discrete logarithm problem.

Keywords: Cryptography, RSA algorithm, Triple DES, Asymmetric key cryptography, 'n' prime number.

Introduction

In the today's era the internet provides communication between people and facilitates for electronic payment, military communication and many others. This cause a major concern for privacy, identify theft, security etc. cryptography is a standard way of secure the data over the medium.

Cryptography has been developed from the Greek word krypton and graphein which means is hiding information person who study and discover cryptography are called cryptographers and study of cryptography is name by cryptanalysis.

Cryptography is a part of secret information. It is science and art of protecting the information over the medium. It is process of convert readable text to unreadable text. By using the cryptography we can help this fickle information by private document on over computer network in a distributed network cryptography become important part of secure communication. There are three type of cryptography algorithm: symmetric key cryptography, Hashing, asymmetric key cryptography.

An algorithm for cryptography that uses the same keys for both encryption of normal text and decryption for cipher text is called symmetric key cryptography. e.g. Data Encryption standard (DES) and Advance Encryption standard (AES). To solve the

key distribution problem Whitfield Diffie and Martin Hellman [11] developed the concept of public key cryptography in 1976. Rivest, Adi Shamir and Leonard Aldeman are discover RSA in 1977. it generates two key: Public key for Encryption and Private Key to Decryption message [3]. RSA algorithm consist of three phase: First phase key generation, Second phase is encryption and Third phase is decryption.

As a public key is work for encryption and is well known to everyone and with the help of public key, defender can use brute force attack method to find corresponding private key which is used to decrypt message and get original message [4].

The proposed algorithm is similar to RSA with few modification. Proposed algorithm is also known by Public key cryptography. In this algorithm we have taken extremely large number that has four prime factor (similar to RSA) in addition of this used to two public

Related work

Cryptography is a process which is associated with encloses plaintext into cipher text (encryption process) then back again plaintext (decryption). In Asymmetric key cryptography using two different keys: Public key and Private Key

.Private key cannot obtain by Public key. This is one major difference between Asymmetric and Symmetric key cryptography, and that major difference change whole process. mostly it has implication throughout the security .As compare symmetric key cryptography as faster move easy and better suited for application drawback of symmetric key cryptography is less secure and move open to wider areas of attacks.

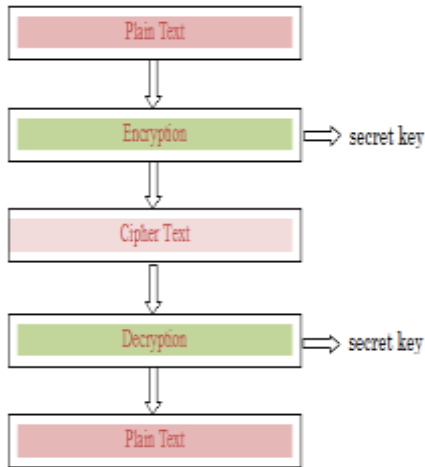


Figure 1 Symmetric Key Cryptography

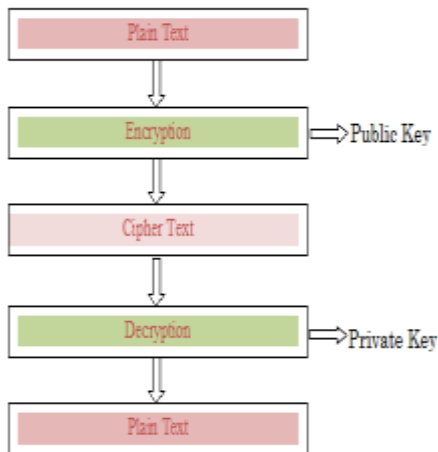


Figure 2 Asymmetric Key Cryptography

Literature review

Rivest, Shamir and Alderman-RSA methodology for confidentially and authentication in this research paper which is used RSA algorithm for secure transmission over the computer and network. It is also increased the efficiency and security. According to Ravi Shankar: Security of RSA algorithm depends on prime number because it is difficult to crack the large prime number. It is

provide the security and performance. In this paper a Modified RSA algorithm is provide security against brute force attack.Hu-Zhou: In this research paper which is used large prime number RSA cryptography for security [3].the prime number is not easily factorized.

A:-The RSA digital signature has appropriate mathematical foundation, which as follow [5]

Theorem 1: Any positive Integer a can be denoted by a_i where

$$A_i = p_1 p_2 p_3 \dots p_n \quad , a_i > 0$$

Theorem 2 :(Euclid theorem): The greatest common divisor g of the positive integer a and b can be represented as a linear sum of original two number a and b .In other world, it is always possible to integer s and t such that –

$$g = s*a + t*b \quad [6]$$

Theorem 3 :(Fermat little theorem): it state that if p is a large prime number, then for any positive integer a, then $a_p = a \text{ mod } (p)$ or $a_{p-1} = 1 \text{ mod } (p)$

Theorem 4: if p and q are prime number and p not equal to q then

$$\phi(p \ q) = \phi(p) * \phi(q) = (p-1) (q-1)$$

B: RSA key generation algorithm-

1. Select two different large random prime number p and q.
2. Calculate $n = p * q$ and $\phi(n) = (p-1) * (q-1)$ [theorem 4]
- Where ϕ is an Euler's function
3. Choose an integer e, such that $1 < e < \phi(n)$ and $\text{gcd}(e, \phi(n)) = 1$ [theorem 2] where e, $\phi(n)$ are co-prime.
4. Compute d:
- d is multiplication inverse of e mod $(\phi(n))$
 $e * d \text{ mod } \phi(n) = 1$
5. The public key is (e, n) and private key is (d, n)

1). Encryption

Sender A know the following

- 1-Recive the receiver B's public key.
- 2-The plaintext message as a positive Integer m.
- 3-calculate the cipher text $C = m^e \text{ mod } n$
- 4-calculated C sends to B

2). Decryption

Receiver B does the following

- 1-Using private key to compute $m = C^d \text{ mod } n$
- 2-Get original plain text m.

Problem domain

The positive large number is easily factorized or break and less prime number are easily decomposed which will not provided more security over the network, that's why we used two public key and n prime number to provide more security over

the network and it is also not easily break and data sharing between different nodes are vulnerable.

and secure data transfer on the transmission medium [12].

Solution methodology

In this research paper we developed an algorithm it is based on modification RSA algorithm ~~SOLUTION METHODOLOGIES~~ prime number. This algorithm provides high security over the network

Proposed solution

A. Process of Encryption

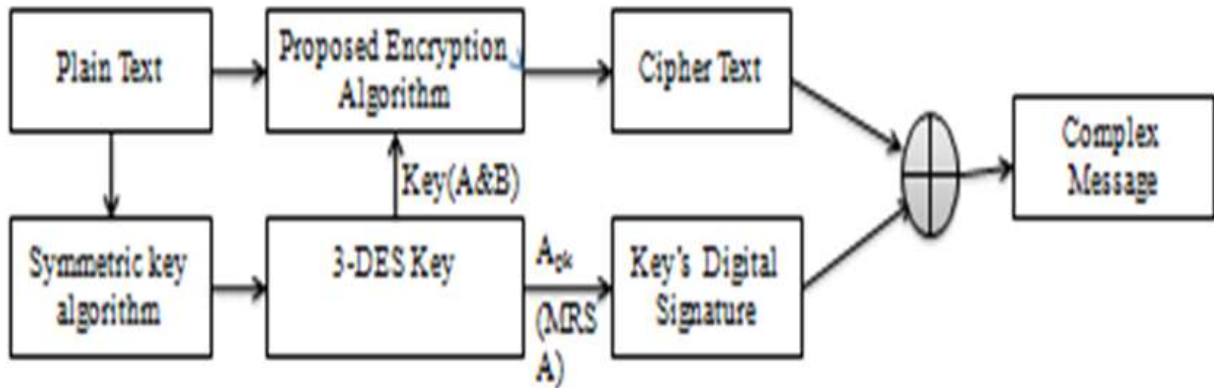


Figure 3 Encryption Process

During the process of sending encrypted information, the random number generator uses 3-DES session key only once, it encrypt the plaintext to output cipher text. On the other hand, the originator get public key from public key management centre, and then using MRSA to encrypt session key. Finally, the mixture of the session key from MRSA encryption and the cipher text from 3-DES encryption are sent out to receiver.

B. Process of Decryption

The decryption of Hybrid proposed algorithm is as follows. Firstly, the receiver B received message and divide received cipher text Complex Message into two parts, first is cipher text key from the MRSA algorithm encryption, and second is cipher text C from the proposed 3-DES algorithm encryption and second, the receiver B decrypt cipher text Key by their own private key, receive the original key K which belongs hybrid algorithm, then decrypt the cipher text C to the original M by key K.

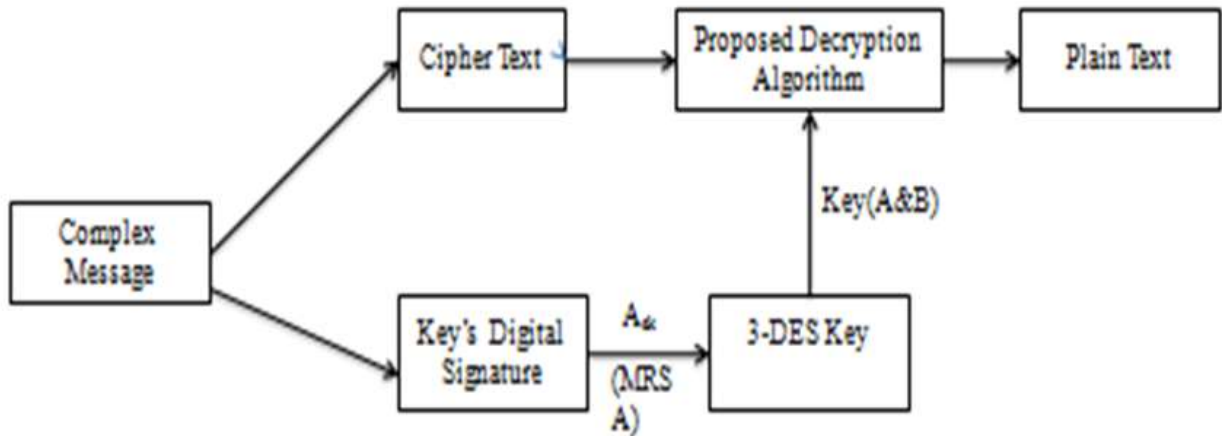


Figure 4 Decryption Process

C. Computational Steps for Key Generation

RSA is 1024 bit block cipher in which the plain text and cipher text integer value lie between 0 to n-1. In which we will be used four prime number and get public key and private key [7] and also using two public keys and one private key for encryption and decryption

- 1). Proposed RSA key generation
 Computational steps for selecting the largest prime number p, q, r and s in RSA cryptography.
 To generated the prime integer p, q, r and s.
 - Firstly, we decided upon the size of integer and implementation if RSA of size B Bits.
 - Using the high quality random number generator [8], you first generated a random number of size B/2 bits
 - We set the lowest bits of integer generated by the above: This insures that the number will be odd [9]
 - We also set the two highest bits of the integer: this insure that the highest bits of null are set.
 - Using the Miller–Rabin-theorem, check to see if the resulting integer is prime .if not you can increment by 2 and check again.

- Calculate $n = p * q * r * s$ and $\phi(n) = (p-1)(q-1)(r-1)(s-1)$
 - Where ϕ is Euler's is function.
 - Choose an integer value e, where e lies between 1 to $\phi(n)$ and $\gcd(e, \phi(n)) = 1$, select two number a and b such that $b = a * e$ and using this number two public key $\{b, e\}$, $\{a\}$ [10].
 - Finally compute d as multiplication inverse of e mod $(\phi(n))$
- 2). Encryption
 Suppose that user A has shared its public key and that user B send the message m to A.
 Then B calculate cipher text $C = (m^{b/a})^d \text{ mod } n$ and then send C.
 - 3). Decryption
 Decryption of the cipher text by A and user A decrypt message $m = C^d \text{ mod } n = (m^{b/a})^d \text{ mod } n = m^{b/ad} \text{ mod } n$.
 Both sender and receiver must known about the values of n, b and a only receiver known the secret value of d In Asymmetric key cryptography with public key $K.U = \{b, n\}, \{a\}$ and private key of $K.R = \{d, n\}$ [10].

Comparison among RSA, modified RSA using Two public key and MRSA using n prime number with n prime number

Table 1 Comparison table among RSA, MRSA and MRSA with n prime number

Serial No.	RSA	Modified RSA	MRSA with n prime number
1.	Use only one public key	Use 2 public key	Use 2 public key
2.	Less communication overhead	Medium communication overhead	High communication overhead
3.	Process speed is fast	Process speed is slow	Process speed is very low
4.	It has less security	It is increasing security	It is provide more security
5.	More permeable to brute force attack	Less permeable to brute force attack	Little permeable to brute force attack
6.	Using encryption and decryption required time is more.	Using encryption and decryption required time is less.	Using encryption and decryption required time is little.

Table 1 shows the general compare among RSA, Modified RSA and MRSA using n prime number .In this algorithm we found that by increasing module length n then increase security and speed decrease. Key generation point of view MRSA, MRSA with n prime number is slower than RSA. In encryption point of view all are working almost same. In case of algorithm only one multiplication operation is additional for each fragment calculation. For decryption point of view MRSA, RSA is almost same .Overall performance vice MRSA with n prime number is better in security but less in speed and throughput.

System overview

With the recent seizure and spread of the data sharing paradigm in distributed systems such as online social networks or cloud computing, there have been increasing demands and concerns for distributed data security. Improving security and efficiency in data sharing over the transmission medium and network [14]

In IESDS three levels of Authentication are provided with a dedicated architecture

A .First Level

It is the User, who is having all the privileges i.e. can add data, can add node, can share to specific node or can sharing to all the workstations members.

B. Second Level

Here only authentication nodes access server

C. Third Level

After second level verification, node access folder created by self then this permission granted by Data Access Control level.

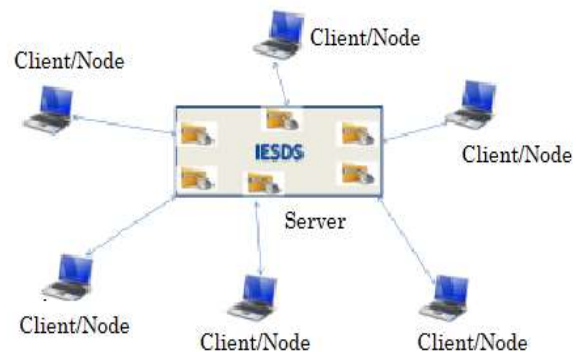


Figure 5. System Overview

Result analysis*Table 2 Plain Text files size with Encryption Time*

File Name	Plain Text (In Kilobyte)	Encryption Time (In Seconds)
Data Set 1	1	0.00001
Data Set 2	4.52	0.00012
Data Set 3	10	0.00039
Data Set 4	50	0.00089
Data Set 5	300	0.00371
Data Set 6	786	0.53125
Data Set 6	1326	1.05123
Data Set 7	5529	4.23467

Table 3 Multimedia files size with Encryption Time

File Name	Plain Text (In Megabyte)	Encryption Time (In Seconds)
Data Set 1	1.87	1.26
Data Set 2	3.52	2.21
Data Set 3	7.39	6.63
Data Set 4	18	13.53
Data Set 5	22.4	16.18
Data Set 6	39.8	31.28
Data Set 6	71	53.21
Data Set 7	99.8	81.34

We apply our algorithm to various size of sample plain text file which the sender wants to transmit through wired or wireless technology and set up the results such as size of text file to be encrypt, size of cipher text file and time taken by plain text file in encryption. Table2 calculate encryption time for

different size of plain text file and multimedia file. In order to compute the effect of change in plain text files in getting the encryption time. We measure that as the size of text file increased the encryption time had also increased.

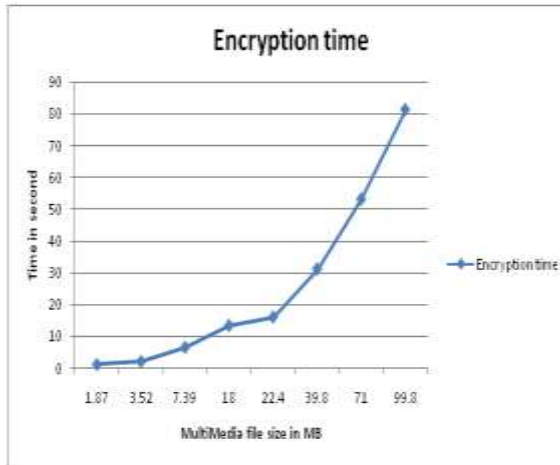


Figure 6 Encryption time for different size of plain text file

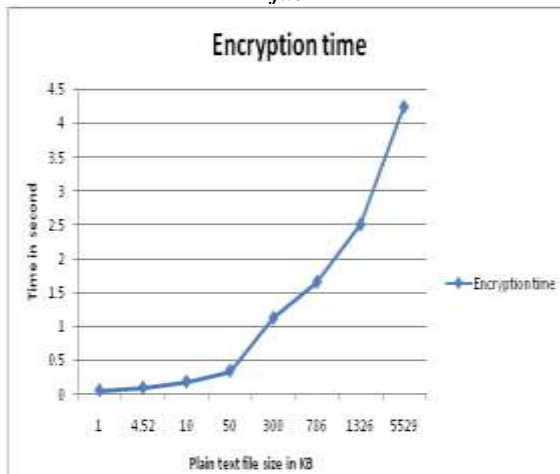


Figure 7 Encryption time for different size of multimedia file

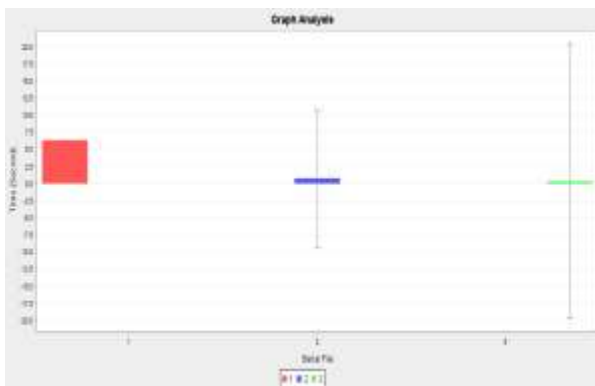


Figure 8 Analysis of plain text size and cipher text size with Encryption time

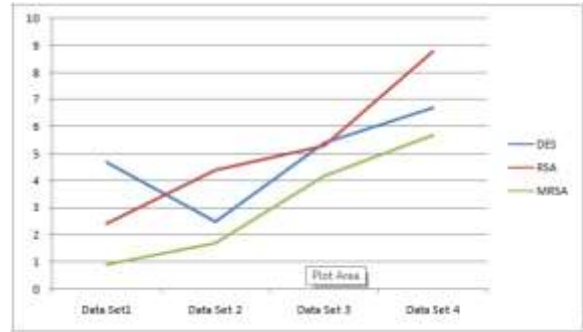


Figure 9 Analyses of DES, RSA and MRSA with Encryption time

At the certain level the time of encryption would grow linearly because of increasing in plain text as shown in fig 6 and fig 7. Table 2 and Table 3 provided time taken by various sizes of plain text file. Our analysis has shown that using of both symmetric and asymmetric key cryptography produce output (cipher text) size is small or integrated as compare to plain text size. This is the advantage of using both the techniques together, which we have achieved from analyzed results. We analyzed that as the file size of plain text is small the size of cipher text is little equal up to some extend and when the file size is greater than 150 KB the size of cipher text is decreased a lot.

Conclusion and future work

In this research paper an algorithm is proposed for RSA a method for implementing a public key cryptography (RSA) using two public and four prime number and same mathematical equations. By using two public keys and n prime number, will provide the security over the network so attacker cannot get keys and unable to decrypt the message. The proposed modified RSA approach is used for system that provides more security but less speed compare to RSA algorithm and improving security and efficiency in data sharing over the network. In our future work we will implement it for advance research such as secure transmission of file, video file, image file, etc. this may perhaps our future research topic using hybrid data encryption and decryption approach.

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